

M269

Programming & Efficiency Topics

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1 Meeting Agenda

- Revue of session on Binary Trees
- Exponentials and Logarithms
- Programming points
- Height balanced (AVL) trees
- Future topics

2 Exponentials and Logarithms

2.1 Exponentials and Logarithms — Definitions

- **Exponential function** $y = a^x$ or $f(x) = a^x$
- $a^n = a \times a \times \cdots \times a$ (n a terms)
- **Logarithm** reverses the operation of exponentiation
- $\log_a y = x$ means $a^x = y$
- $\log_a 1 = 0$
- $\log_a a = 1$
- Method of logarithms propounded by John Napier from 1614
- **Log Tables** from 1617 by Henry Briggs
- **Slide Rule** from about 1620–1630 by William Oughtred of Cambridge
- *Logarithm* from Greek *logos* ratio, and *arithmos* number (Chambers Dictionary 13th edition 2014)

2.2 Rules of Indices

$$1. a^m \times a^n = a^{m+n}$$

$$2. a^m \div a^n = a^{m-n}$$

$$3. a^{-m} = \frac{1}{a^m}$$

$$4. a^{\frac{1}{m}} = \sqrt[m]{a}$$

$$5. (a^m)^n = a^{mn}$$

$$6. a^{\frac{n}{m}} = \sqrt[m]{a^n}$$

$$7. a^0 = 1 \text{ where } a \neq 0$$

- **Exercise** Justify the above rules
- What should 0^0 evaluate to ?
- See [Wikipedia: Exponentiation](#)
- The *justification* above probably only worked for whole or *rational* numbers — see later for exponents with *real* numbers (and the value of *logarithms, calculus...*)

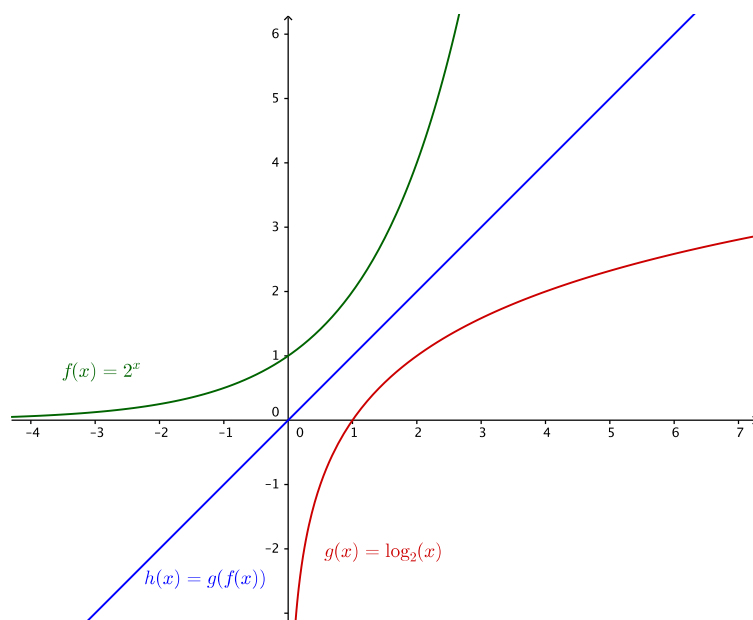
2.3 Logarithms — Motivation

- Make arithmetic easier — turns multiplication and division into addition and subtraction (see later)
- Complete the range of **elementary functions** for differentiation and integration

- An **elementary function** is a function of one variable which is the composition of a finite number of arithmetic operations ((+), (-), ($\times$), ($\div$)), exponentials, logarithms, constants, and solutions of algebraic equations (a generalization of nth roots).
- The elementary functions include the trigonometric and hyperbolic functions and their inverses, as they are expressible with complex exponentials and logarithms.
- See A Level FP2 for Euler's relation $e^{i\theta} = \cos \theta + i \sin \theta$
- In A Level C3, C4 we get $\int \frac{1}{x} = \log_e |x| + C$
- e is **Euler's number** 2.71828...

2.4 Exponentials and Logarithms — Graphs

- See GeoGebra file [expLog.ggb](#)



2.5 Laws of Logarithms

- **Multiplication law** $\log_a xy = \log_a x + \log_a y$
- **Division law** $\log_a \left(\frac{x}{y}\right) = \log_a x - \log_a y$
- **Power law** $\log_a x^k = k \log_a x$

- **Proof of Multiplication Law**

$$x = a^{\log_a x}$$

$$y = a^{\log_a y}$$

by definition of log

$$xy = a^{\log_a x} a^{\log_a y}$$

$$= a^{\log_a x + \log_a y}$$

by laws of indices

$$\text{Hence } \log_a xy = \log_a x + \log_a y$$

by definition of log

2.6 Arithmetic and Inverses

- *Notation helps or maybe not ?*
- **Addition** $\text{add}(b, x) = x + b$
- **Subtraction** $\text{sub}(b, x) = x - b$
- Inverse $\text{sub}(b, \text{add}(b, x)) = (x + b) - b = x$
- **Multiplication** $\text{mul}(b, x) = x \times b$
- **Division** $\text{div}(b, x) = x \div b = \frac{x}{b} = x/b$
- Inverse $\text{div}(b, \text{mul}(b, x)) = (x \times b) \div b = \frac{(x \times b)}{b} = x$
- **Exponentiation** $\text{exp}(b, x) = b^x$
- **Logarithm** $\text{log}(b, x) = \log_b x$
- Inverse $\text{log}(b, \text{exp}(b, x)) = \log_b(b^x) = x$
- *What properties do the operations have that work (or not) with the notation ?*

Arithmetic Operations — Commutativity and Associativity

- **Commutativity** $x \oplus y = y \oplus x$
- **Associativity** $(x \oplus y) \oplus z = x \oplus (y \oplus z)$
- (+) and (×) are *semantically* commutative and associative — so we can leave the brackets out
- (−) and (÷) are not
- Evaluate $(3 - (2 - 1))$ and $((3 - 2) - 1)$
- Evaluate $(3/(2/2))$ and $((3/2)/2)$
- We have the *syntactic* ideas of left (and right) associativity
- We choose (−) and (÷) to be left associative
- $3 - 2 - 1$ means $((3 - 2) - 1)$
- $3/2/2$ means $((3/2)/2)$
- **Operator precedence** is also a choice (remember BIDMAS or BODMAS ?)
- If in doubt, put the brackets in

Exponentials and Logarithm — Associativity

- What should 2^{3^4} mean ?
- Let $b^x \equiv b^x$
- Evaluate $(2^3)^4$ and $2^{(3^4)}$
- Evaluate $c = \log_b(\log_b((b^b)^x))$
- Evaluate $d = \log_b(\log_b(b^{(b^x)}))$
- Beware spreadsheets Excel and LibreOffice here
- $(2^3)^4 = 2^{12}$ and $2^{3^4} = 2^{81}$
- Exponentiation is not semantically associative
- We *choose* the syntactic left or right associativity to make the syntax nicer.
- Evaluate $c = \log_b(\log_b((b^b)^x))$
- $c = \log_b(x \log_b(b^b)) = \log_b(x \cdot (b \log_b b)) = \log_b(x \cdot b \cdot 1)$
- Hence $c = \log_b x + \log_b b = \log_b x + 1$
- Not symmetrical (unless b and x are both 2)
- Evaluate $d = \log_b(\log_b(b^{(b^x)}))$
- $d = \log_b((b^x)(\log_b b)) = \log_b((b^x) \times 1)$
- Hence $d = \log_b(b^x) = x(\log_b b) = x$
- Which is what we want — so exponentiation is *chosen* to be right associative

2.7 Change of Base

- Change of base

$$\log_a x = \frac{\log_b x}{\log_b a}$$

Proof: Let $y = \log_a x$

$$a^y = x$$

$$\log_b a^y = \log_b x$$

$$y \log_b a = \log_b x$$

$$y = \frac{\log_b x}{\log_b a}$$

- Given x , $\log_b x$, find the base b

$$b = x^{\frac{1}{\log_b x}}$$

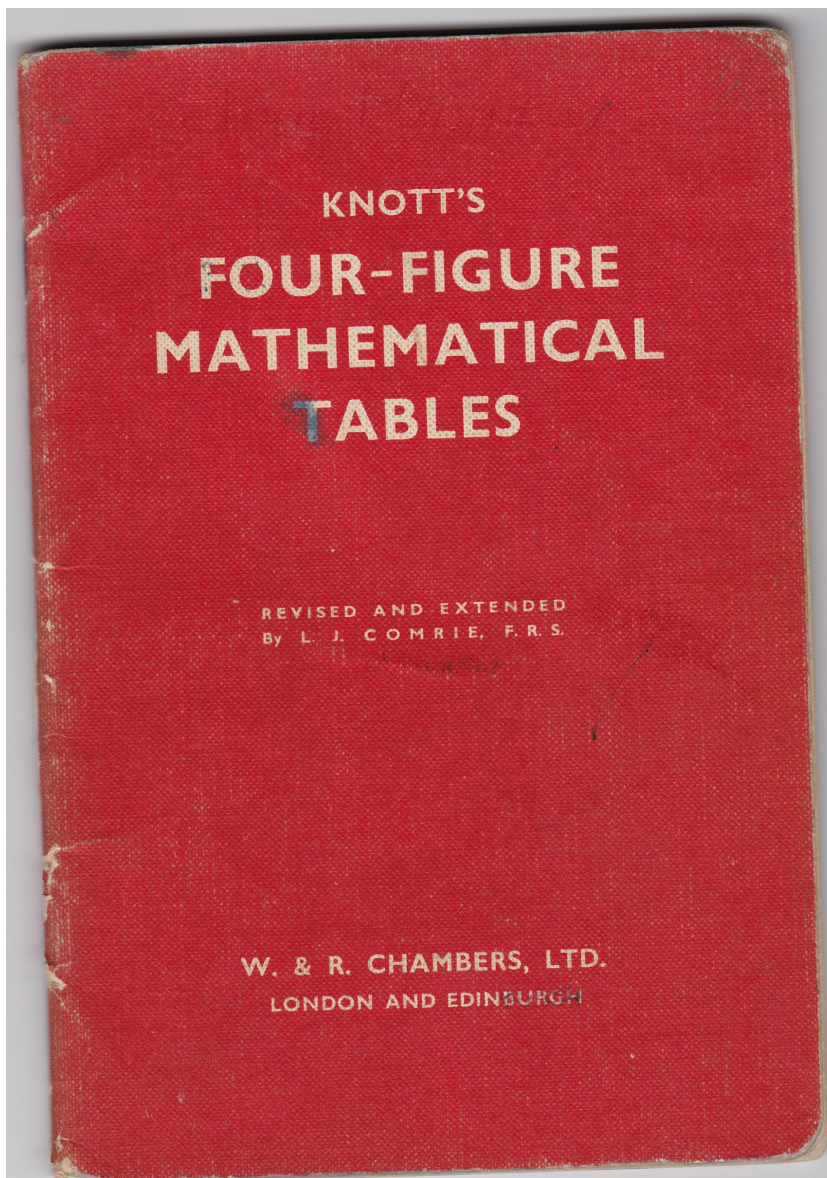
- $\log_a b = \frac{1}{\log_b a}$

3 Before Calculators and Computers

- We had computers before 1950 — they were *humans* with pencil, paper and some further aids:
- **Slide rule** invented by [William Oughtred](#) in the 1620s — major calculating tool until [pocket calculators](#) in 1970s
- **Log tables** in use from early 1600s — method of logarithms propounded by [John Napier](#)
- **Logarithm** from Greek *logos* ratio, and *arithmos* number

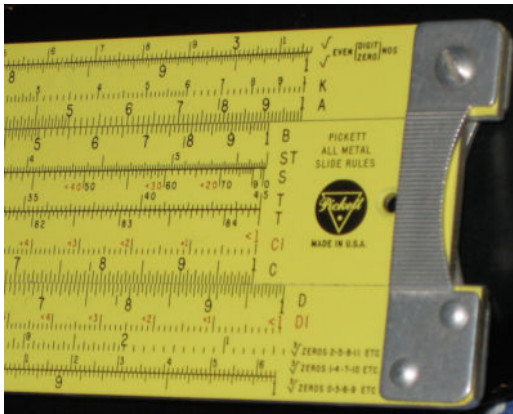
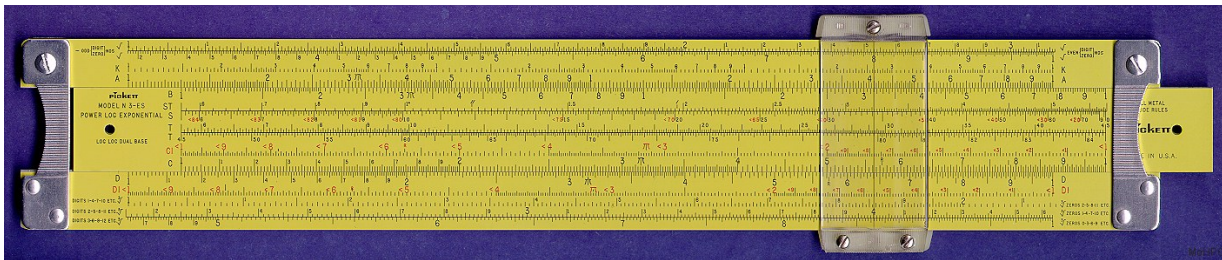
3.1 Log Tables

Knott's Four-Figure Mathematical Tables

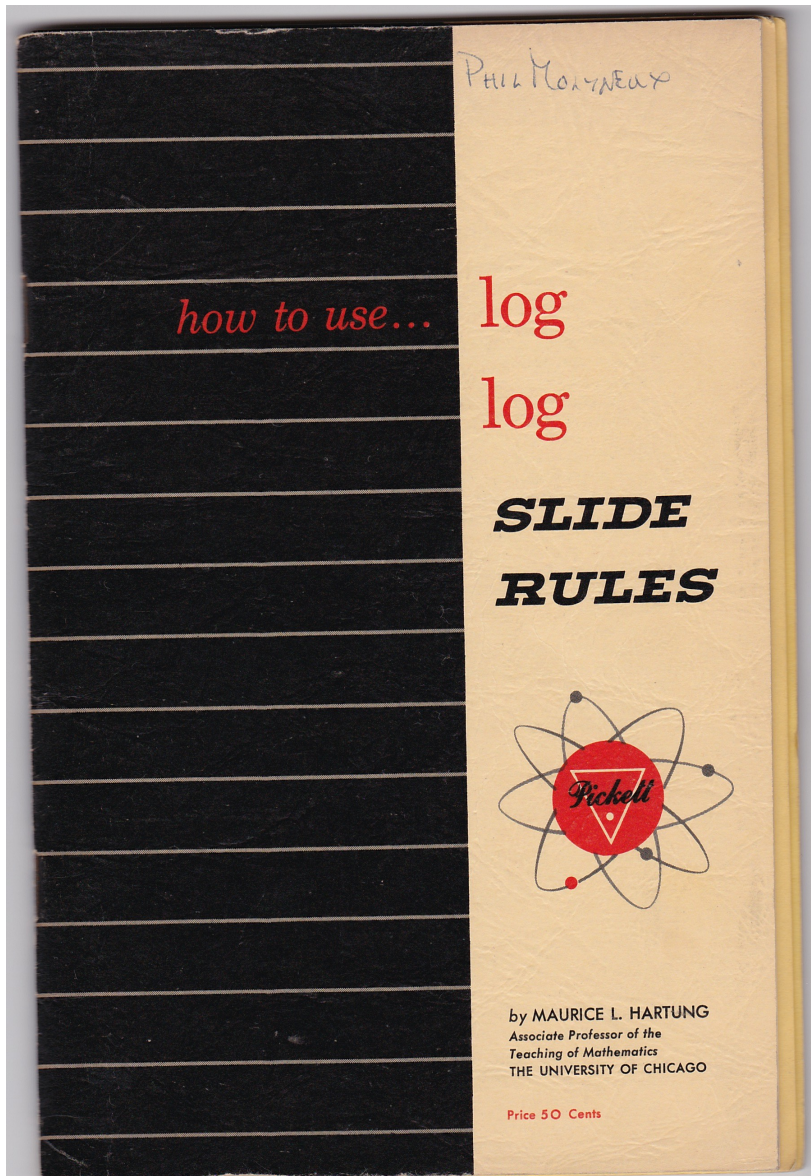


3.2 Slide Rules

Pickett N 3-ES from 1967



- See [Oughtred Society](#)
- [UKSRC](#)
- [Rod Lovett's Slide Rules](#)
- [Slide Rule Museum](#)

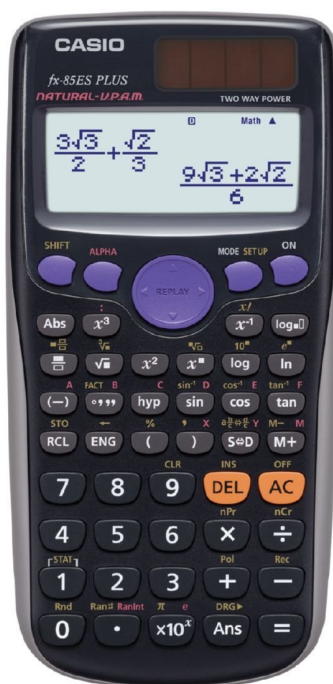
Pickett log log Slide Rules Manual 1953

3.3 Calculators

HP HP-21 Calculator from 1975 £69



Casio fx-85GT PLUS Calculator from 2013 £10



Calculator links

- HP Calculator Museum <http://www.hpmuseum.org>
- HP Calculator Emulators <http://nonpareil.brouhaha.com>
- HP Calculator Emulators for OS X <http://www.bartosiak.org/nonpareil/>
- Vintage Calculators Web Museum <http://www.vintagecalculators.com>

3.4 Example Calculation

- Evaluate 89.7×597
- **Knott's Tables**
- $\log_{10} 89.7 = 1.9528$ and $\log_{10} 597 = 2.7760$
- Shows *mantissa* (decimal) & *characteristic* (integral)
- Add 4.7288, take *antilog* to get $5346 + 10 = 5.356 \times 10^4$
- **HP-21 Calculator** — set display to 4 decimal places
- $89.7 \text{ [log]} = 1.9528$ and $597 \text{ [log]} = 2.7760$
- $+$ displays 4.7288
- $10 \text{ [ENTER]}, x \neq y$ and y^x displays 53550.9000
- **Casio fx-85GT PLUS**
- $\text{[log]} 89.7 \text{ [)]} = 1.952792443 \text{ [+] } \text{[log]} 597 \text{ [)]} = 2.775974331 \text{ [=]}$
- $4.728766774 \text{ [Ans]} + \text{[10^x]}$ gives 53550.9

4 Basic Computational Components

Computational Components — Imperative

Imperative or procedural programming has statements which can manipulate global memory — statements can be organised into procedures (or functions)

- **Sequence** of statements

```
stmt ; stmt
```

- **Iteration** to repeat statements

```
while expr :  
  suite  
  
for targetList in exprList :  
  suite
```

- **Selection** choosing between statements

```
if expr : suite  
elif expr : suite  
else : suite
```

Functional programming treats computation as the evaluation of expressions and the definition of functions (in the mathematical sense)

- **Function composition** to combine the application of two or more functions — like sequence but from right to left (notation accident of history)

```
(f . g) x = f (g x)
```

- **Recursion** — function definition defined in terms of calls to itself (with *smaller* arguments) and base case(s) which do not call itself.
- **Conditional expressions** choosing between alternatives expressions

```
if expr then expr else expr
```

4.1 Writing Programs & Thinking

The Steps

1. Invent a *name* for the program (or function)
2. What is the *type* of the function ? What sort of *input* does it take and what sort of *output* does it produce ?
3. Invent *names* for the input(s) to the function (*formal parameters*)
4. *Think* of the definition of the function body.

The *Think* Step

0. Think of an example or two — what should the program/function do ?
1. Don't think too much at one go — break the problem down. Top down design, step-wise refinement.
2. What are the inputs — describe all the cases.
3. Investigate choices. What data structures ? What algorithms ?
4. Use common tools — bottom up synthesis.
5. Spot common function application patterns — generalise & then specialise.
6. Look for good *glue* — to combine little programs to make bigger ones.
7. Try out your first examples when you have written the program

5 Divide and Conquer

Binary Search — Exercise

Given the *Python* definition of `binarySearchRec` from the slides for the Unit 1 tutorial, trace an evaluation of `binarySearchRec(xs, 25)` where `xs` is

```
xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
```

Binary Search Recursive

```

1  def binarySearchRec(xs, val, lo=0, hi=-1):
2      if (hi == -1):
3          hi = len(xs) - 1
4
5      mid = (lo + hi) // 2
6
7      if hi < lo:
8          return None
9      else:
10         guess = xs[mid]
11         if val == guess:
12             return mid
13         elif val < guess:
14             return binarySearchRec(xs, val, lo, mid-1)
15         else:
16             return binarySearchRec(xs, val, mid+1, hi)

```

Binary Search — Solution

```

xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
binarySearchRec(xs, 25)
xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
binarySearchRec(xs,25,8,14) by line 15
xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
binarySearchRec(xs,25,8,10) by line 15
xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
binarySearchRec(xs,25,8,8) by line 13
xs = [2,5,7,15,17,19,21,24,27,31,37,48,57,87,95]
binarySearchRec(xs,25,8,7) by line 13
Return value: None by line 7

```

6 String Search

- *Sunday Quick Search* algorithm
- *Knuth-Morris-Pratt (KMP)* algorithm
- [Exact String Matching Algorithms](#) — animations in Java

6.1 Sunday Quick Search Algorithm

Example

- Sunday Quick Search example from M269 2014J TMA02
- Consider the alphabet C, G, A, T and a target string CGTACTCGTAGT.
- Calculate the shift table for this search, explaining in detail how you used the part of the algorithm that builds the table to derive your results.
- Given the target string CGTACTCGTAGT, the search string CGTACTCGGCGTAAAGT-CCGTCTT and the shift table calculated above

- describe, with diagrams if necessary, how the first attempt to match the target string against the search string fails
- and how the target string slides along the search string for the second matching attempt.

Sunday Quick Search Shift Table

- The shift table for the Sunday Quick Search algorithm:
 - If the character does not appear in the target string T , the shift distance is one more than the length of T
 - If the character does appear in T the shift distance is the first position at which it appears, counting from right to left and starting at 1
- Shift table:

A	C	G	T
3	6	2	1
- See the example at [Quick Search algorithm](#)

Calculating the Shift

- Given the following alignment of *Search* and *Target* strings

Search string

C	G	T	A	C	T	C	G	G	C	G	T	A	A	A	G	T	G	C	G	T	C	T	T
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

C	G	T	A	C	T	C	G	T	A	G	T
---	---	---	---	---	---	---	---	---	---	---	---

Target string

- The next character in the *Search* string after the *Target* string is **A** so the shift will be 3 places

Search string

C	G	T	A	C	T	C	G	G	C	G	T	A	A	A	G	T	G	C	G	T	C	T	T
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

C	G	T	A	C	T	C	G	T	A	G	T
---	---	---	---	---	---	---	---	---	---	---	---

Target string

6.2 Knuth-Morris-Pratt (KMP) Algorithm

Example

- Knuth-Morris-Pratt (KMP) algorithm example from M269 2014J TMA02
- Consider the alphabet C, G, A, T and a target string CGTACTCGTAGT.
- Calculate the prefix table for this target string,

- explaining in detail how you derived the sixth, seventh and eighth entries in your table.

KMP Prefix Table

- The KMP *prefix table* takes the *target string*, t , and a *position*, p , in t and returns an integer k
- k is the length of the maximum proper prefix of t up to position p that is a suffix of t up to position p
- We define a prefix function to take a *target string*, t , and a number $k \geq 0$, which is the number of items off the front of t
- Formally, if our position index is 0 based we have:
- $\text{prefixTable}(t, 0) = 0$
- $\text{prefixTable}(t, p) = \max\{k : k \leq p \wedge \text{prefix}(t, k) \sqsubset \text{prefix}(t, p + 1)\}$
- Here $\sqsubset$ means *proper suffix* — that is a suffix that does not include the whole string

KMP Shift Function

- M269 does not give the KMP shift function (or table) — we give it here for interest
- The shift function takes the *target string*, t , and the number of characters in the target string that have been matched, q and returns the shift.
- $\text{shift}(t, 0) = 1$
- $\text{shift}(t, q) = q - \text{prefixTable}(t, q - 1)$
- This assumes 0 based list indexing
- The notation here comes from [Cormen et al. \(2009, 2009, section 32.4, page 1004\)](#) — this uses 1 based list indexing and also provides code.
- Cormen uses the terminology *Text* for *Search string* and *Pattern* for *Target string*

KMP Prefix and Shift Table

C	G	T	A	C	T	C	G	T	A	G	T	Target string
0	1	2	3	4	5	6	7	8	9	10	11	Position (Index), p
0	1	2	3	4	5	6	7	8	9	10	11	Match, q
0	0	0	0	1	0	1	2	3	4	0	0	$\text{prefixTable}(t, p)$
1	1	2	3	4	4	6	6	6	6	6	11	$\text{shift}(t, q)$

7 Future Work

- Reflections on today's topics — what were the important points ?
- Dates for next meeting: (??)
- Next topics
 - Unit 5 Optimisation
 - Graph algorithms
 - Critique of Phil's notes

8 Web Sites & References

8.1 Web Sites

- *Python Documentation* <https://docs.python.org/3/> (26 January 2016)
- *Wikipedia Category: Python Libraries* https://en.wikipedia.org/wiki/Category:Python_Libraries (26 January 2016)
- *Python Wiki: Useful Modules* <https://wiki.python.org/moin/UsefulModules> (26 January 2016)
- *Python Packaging User Guide* <http://python-packaging-user-guide.readthedocs.org/> (26 January 2016)
- *Python Packaging Authority* <https://www.pypa.io/en/latest/> (26 January 2016)
- *PyPI Python Package Index* <https://pypi.python.org/pypi> (26 January 2016)
- *Top 100 Tools for Learning 2015* <http://c4lpt.co.uk/directory/top-100-tools/> (26 January 2016)
- *Pygal Python charting* <http://www.pygal.org/en/latest/> (26 January 2016)
- *21 Ridiculously Impressive HTML5 Canvas Experiments* <http://code.tutsplus.com/articles/21-ridiculously-impressive-html5-canvas-experiments--net-14210> (26 January 2016)
- *graph-tool Efficient network analysis* <https://graph-tool.skewed.de> (26 January 2016)
- *Python Patterns — Implementing Graphs* <https://www.python.org/doc/essays/graphs/> (26 January 2016)
- *Graphs in Python* http://www.python-course.eu/graphs_python.php
- *Stackoverflow Python Graph Library* <http://stackoverflow.com/questions/606516/python-graph-library> (26 January 2016)
- *Dijkstra's algorithm for shortest paths* <http://code.activestate.com/recipes/119466-dijkstras-algorithm-for-shortest-paths/>

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